

Introduction to graph theory

Graphs

- Size and order
- Degree and degree distribution
- Subgraphs
- Paths, components
- Geodesics
- Some special graphs
- Centrality and centralisation

Directed graphs

- Dyad and triad census
- Paths, semipaths, geodesics, strong and weak components
- Centrality for directed graphs
- Some special directed graphs

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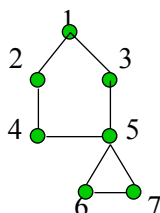
Definition of a graph

A **graph** G comprises a set V of vertices and a set E of edges

Each **edge** in E is a pair (a,b) of vertices in V

If (a,b) is an edge in E , we connect a and b in the **graph drawing** of G

Example:



$V = \{1, 2, 3, 4, 5, 6, 7\}$
 $E = \{(1, 2), (1, 3), (2, 4),$
 $(4, 5), (3, 5), (4, 5),$
 $(5, 6), (6, 7)\}$

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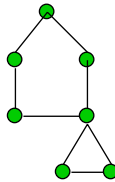
Size and order

The *size* of G is the number n of vertices in V

The *order* of G is the number L of edges in E

Minimum possible order is 0 (*empty* graph)

Maximum possible order is $n(n-1)/2$ (*complete* graph)



Size = 7, Order = 8

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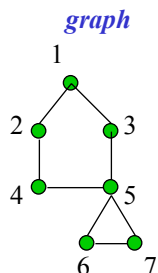
Adjacency matrix for a graph

The adjacency matrix $\mathbf{x} = [x_{ab}]$ for G is a matrix with n rows and n columns and entries given by:

$x_{ab} = 1$ if (a,b) is an edge in G

0 otherwise

Example:



adjacency matrix

```
0110000
1001000
1000100
0100100
0011011
0000101
0000110
```

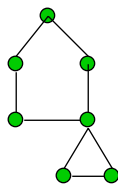
symmetric

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Density

The **density** of G is the ratio of edges in G to the maximum possible number of edges

$$\text{Density} = \frac{2L}{n(n-1)}$$



$$\text{Density} = \frac{2 \times 8}{(7 \times 6)} = \frac{8}{21}$$

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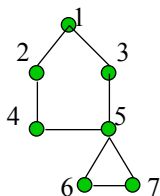
Degrees and degree sequence

The **degree** d_a of vertex a is the number of vertices to which a is linked by an edge

The minimum possible degree is 0

The maximum possible degree is $n-1$

The **degree sequence** for a graph is the vector (d_1, d_2, \dots, d_n)

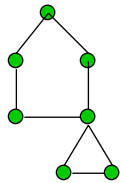


$$\text{Degree sequence} = (2, 2, 2, 2, 4, 2, 2)$$

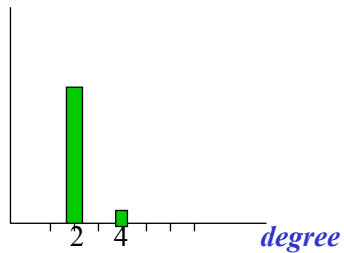
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Degree distribution

The *degree distribution* for the graph is $(k_0, k_1, \dots, k_{n-1})$,
where k_j = the number of nodes with degree j



frequency

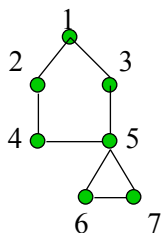


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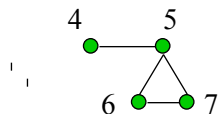
Subgraphs

A *subgraph* of $G=G(V,E)$ is a subset W of the vertex set V together with all of the edges that connect pairs of vertices in W

Eg if $W=\{4,5,6,7\}$, the subgraph of



is



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Subgraph counts: the dyad census

The graph G has $n(n-1)/2$ subgraphs of size 2

Each subgraph of size 2 comprises a pair of vertices, and the edge between them is either present or absent:

subgraph

count



$$D_0 = n(n-1)/2 - L$$



$$D_1 = L$$

Dyad census = (D_0, D_1) *count of the no. of each type of dyad subgraph*

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Subgraph counts: the triad census

The graph G has $n(n-1)(n-2)/6$ subgraphs of size 3

Each subgraph of size 3 comprises a triple of vertices, and the possible forms are:

subgraph count



$$T_0 = (1/6) \sum_{i,j,k} (1-x_{ij})(1-x_{ik})(1-x_{jk})$$



$$T_1 = (1/6) \sum_{i,j,k} (1-x_{ij})(1-x_{ik})x_{jk}$$



$$T_2 = (1/6) \sum_{i,j,k} (1-x_{ij})x_{ik}x_{jk}$$



$$T_3 = (1/6) \sum_{i,j,k} x_{ij}x_{ik}x_{jk}$$



Triad census = (T_0, T_1, T_2, T_3)

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Geodesics

A **geodesic** from a to b is a path of minimum length

The **geodesic distance** d_{ab} between a and b is the length of the geodesic

If there is no path from a to b , the geodesic distance is **infinite**

For the graph

The geodesic distances are:

$$d_{AB} = 1, d_{AC} = 1, d_{AD} = 1, d_{BC} = 1, d_{BD} = 2, d_{CD} = 2$$

Cycles

A **cycle** is an ordered sequence

$$a=v_0, v_1, \dots, v_m=a$$

of vertices in which each adjacent pair (v_{j-1}, v_j) of vertices is linked by an edge, and v_0, v_1, \dots, v_{m-1} are distinct. The **length** of the cycle is m

Cycles of length

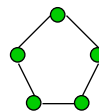
3



4

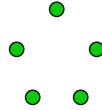


5

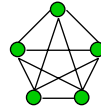


Some special graphs: trivial, empty and complete graphs

The *empty* graph on 5 vertices (Z_5)



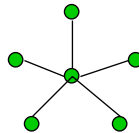
The *complete* graph on 5 vertices (K_5)



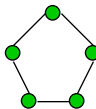
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Star and cyclic graphs

A *star* graph on 6 vertices



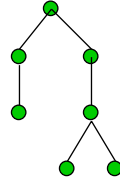
A *cyclic* graph on 5 vertices (C_5)



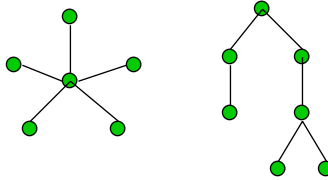
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Trees and forests

A **tree** (a connected acyclic graph)



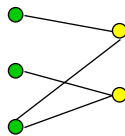
A **forest** (a graph with tree components)



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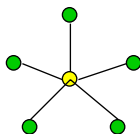
Bipartite graphs

A **bipartite** graph (vertex set can be partitioned into 2 subsets, and there are no edges linking vertices in the same set)

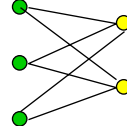


A **complete** bipartite graph (all possible edges are present)

$K_{1,5}$



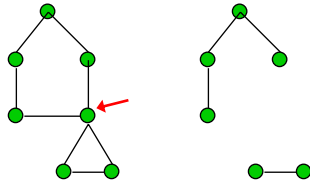
$K_{3,2}$



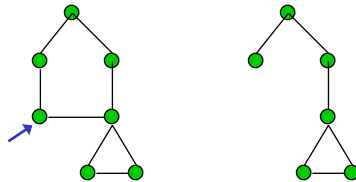
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Cutpoints

A vertex is a **cutpoint** if its removal increases the number of components in the graph



the vertex marked by the red arrow is a cutpoint

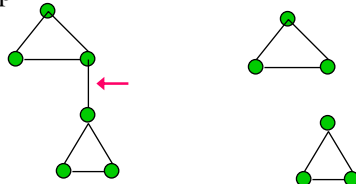


the vertex marked by the blue arrow is not

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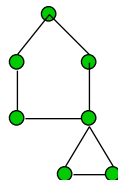
Bridges

An edge is a **bridge** if its removal increases the number of components in the graph



the edge marked by the red arrow is a bridge

This graph has no bridges



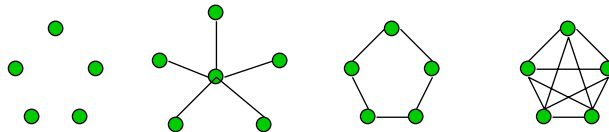
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Connectivity

The **connectivity** $\kappa(G)$ of a connected graph G is the minimum number of vertices that need to be removed to disconnect the graph (or make it empty)

A graph with more than one component has connectivity 0

Graph



Connectivity

0 1 2 4

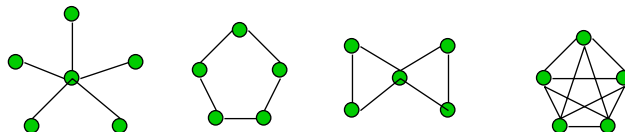
A graph with connectivity k is termed ***k-connected***

Edge-connectivity

The **edge-connectivity** $\lambda(G)$ of a connected graph G is the minimum number of **edges** that need to be removed to disconnect the graph

A graph with more than one component has edge-connectivity 0

Graph



Edge-Connectivity

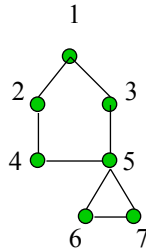
1 2 2 4

Connectivity

1 2 1 4

Independent and edge-independent paths

Two paths from a to b are **independent** if they have no nodes in common apart from a and b e.g. paths $1-2-4-5$ and $1-3-5$



Two paths from a to b are **edge-independent** if they have no edges in common e.g. paths $1-2-4-5-6$ and $1-3-5-7-6$

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Several theorems about connectivity

Whitney's theorem

For any graph G , $\kappa(G) \leq \lambda(G) \leq \delta(G)$, where $\delta(G)$ is the minimum degree of any vertex in G

Menger's theorem

A graph G is k -connected if and only if any pair of vertices in G are linked by at least k independent paths

Menger's theorem

A graph G is k -edge-connected if and only if any pair of vertices in G are linked by at least k edge-independent paths

For application, see Harary & White (2001)

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Degree Centrality

Freeman (1979) described three measures of *vertex* centrality:

Degree centrality (communication potential)

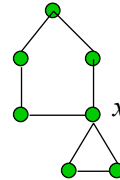
Degree centrality of node a : $C_D(a) = d_a$ **degree of node a**

Normalised degree centrality of node a : $d_a/(n-1)$

Example

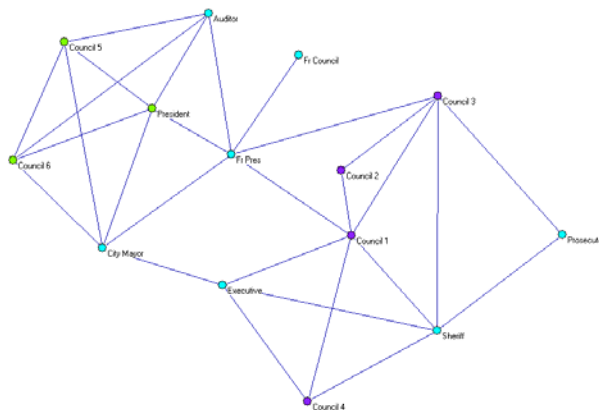
Node x : degree centrality = 4

normalised degree centrality = $4/6 = 0.67$



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A political network (Doreian, 1988)



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Degree centrality in Doreian's (1988) political network

	Degree	NrmDegree	Share
4 Council 1	6.000	46.154	0.107
12 Fr Pres	6.000	46.154	0.107
3 Sheriff	5.000	38.462	0.089
8 President	5.000	38.462	0.089
6 Council 3	5.000	38.462	0.089
13 City Mayor	5.000	38.462	0.089
2 Auditor	4.000	30.769	0.071
1 Executive	4.000	30.769	0.071
9 Council 5	4.000	30.769	0.071
10 Council 6	4.000	30.769	0.071
7 Council 4	3.000	23.077	0.054
14 Prosecutor	2.000	15.385	0.036
5 Council 2	2.000	15.385	0.036
11 Fr Council	1.000	7.692	0.018

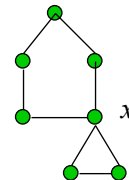
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Closeness Centrality

Closeness centrality (potential for independent communication)

Closeness centrality of node a : $C_D(a) = 1/\sum_b d_{ab}$ **inverse sum of distances to other nodes b**

Normalised closeness centrality of node a : $(n-1)/\sum_b d_{ab}$



Example

Node x : closeness centrality = $1/[1+1+1+1+2+2]=1/8 = 0.125$
 normalised closeness centrality = $6/8=0.75$

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Closeness centrality in Doreian's political network

		Farness	nCloseness
12	Fr Pres	20.000	65.000
4	Council 1	22.000	59.091
6	Council 3	23.000	56.522
13	City Mayor	23.000	56.522
1	Executive	24.000	54.167
8	President	25.000	52.000
3	Sheriff	26.000	50.000
2	Auditor	27.000	48.148
7	Council 4	28.000	46.429
9	Council 5	31.000	41.935
10	Council 6	31.000	41.935
11	Fr Council	32.000	40.625
5	Council 2	32.000	40.625
14	Prosecutor	32.000	40.625

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Betweenness centrality

Betweenness centrality (Potential for control of communication)

Betweenness centrality of node a : $C_D(a) = \sum_{b < c} [g_{bc(a)} / g_{bc}]$

Where g_{bc} is the number of geodesics between b and c , and $g_{bc(a)}$ is the number of geodesics between b and c that contain a

sum over all pairs (b,c) of the proportion of geodesics linking the pair that contain node a

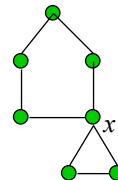
Normalised betweenness centrality of node a :

$$2 \sum_{b < c} [g_{bc(a)} / g_{bc}] / [n^2 - 3n + 2]$$

Example

Node x : betweenness centrality = 14

$$\text{normalised betweenness centrality} = 14 / [49 - 21 + 2] = 7/15$$



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Betweenness centrality in Doreian's political network

Betweenness centrality

	<i>33.198</i>	<i>42.561</i>
12 Fr Pres		
13 City Mayor	14.490	18.578
4 Council 1	13.843	17.747
6 Council 3	13.024	16.697
1 Executive	9.452	12.118
3 Sheriff	4.500	5.769
8 President	4.021	5.156
2 Auditor	3.571	4.579
9 Council 5	0.450	0.577
10 Council 6	0.450	0.577
11 Fr Council	0.000	0.000
5 Council 2	0.000	0.000
7 Council 4	0.000	0.000
14 Prosecutor	0.000	0.000

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Maximum centrality

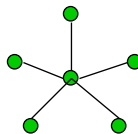
Centrality is at a maximum on all measures for the *central node in a star* configuration:

Normalised measures

Degree centrality: 1

Closeness centrality: 1

Betweenness centrality: 1



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Centralisation

Graph-level measure of *centralisation*:

Degree to which the centrality of the most central vertex exceeds the centrality of all other vertices, compared to the maximum possible discrepancy

Index has the form:

Sum over nodes a of (max centrality in G – centrality of node a)

Max value of the sum in all graphs on the same number of vertices

Can be used with any centrality measure

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Eigenvector centrality (Bonacich, 1972)

The *eigenvector centrality* of vertex a is the sum of its connections to other nodes, weighted by *their* centrality

It is hence given by the solution for c_a of the equation

$$c_a = (1/\lambda) \sum_b x_{ab} c_b \quad \text{where } \lambda \text{ is a constant}$$

[*Mathematical note*: this is equivalent to:

$$\lambda \mathbf{c} = \mathbf{x} \mathbf{c}$$

where \mathbf{x} is the adjacency matrix and \mathbf{c} is the vector of centrality measures; hence \mathbf{c} is an eigenvector of \mathbf{x} , usually taken to be the one associated with the largest eigenvalue λ]

Use with connected graphs only

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Eigenvector centrality for Doreian network

		Eigenvec	nEigenvec
		-----	-----
1	Executive	0.249	35.207
2	Auditor	0.287	40.585
3	Sheriff	0.257	36.368
4	Council 1	0.328	46.358
5	Council 2	0.134	18.887
6	Council 3	0.270	38.154
7	Council 4	0.186	26.357
8	President	0.350	49.435
9	Council 5	0.282	39.854
10	Council 6	0.282	39.854
11	Fr Council	0.083	11.723
12	Fr Pres	0.371	52.454
13	City Mayor	0.343	48.453
14	Prosecutor	0.118	16.655

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Directed graph

A **directed graph** G comprises a set V of vertices and a set E of **arcs**

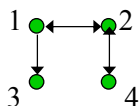
Each **arc** in E is an **ordered pair** (a,b) of vertices in V

If (a,b) is an arc in E , we draw an arc **from a to b** in the **graph drawing** of G

Example:

$$V = \{1, 2, 3, 4\}$$

$$E = \{(1,2), (2,1), (2,4), (1,3), (4,2)\}$$



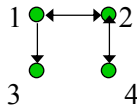
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Adjacency matrix for a directed graph

The adjacency matrix $x = [x_{ab}]$ for G is a matrix with n rows and n columns and entries given by:

$$x_{ab} = \begin{cases} 1 & \text{if } (a,b) \text{ is an arc in } G \\ 0 & \text{otherwise} \end{cases}$$

Example



adjacency matrix

0110

1001

0000

0100

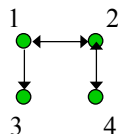
*not
necessarily
symmetric*

Density for a directed graph

The **density** of G is the ratio of arcs in G to the maximum possible number of arcs

$$\text{Density} = \frac{L}{n(n-1)}$$

Example



density = 5/12

Indegrees and outdegrees

The **indegree** i_a of vertex a is the number of vertices to which a is linked by an arc

The **outdegree** o_a of vertex a is the number of vertices linked to a by an arc

The minimum possible indegree (or outdegree) is 0

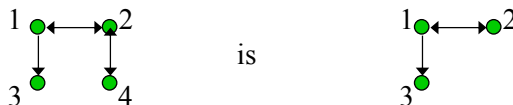
The maximum possible indegree (or outdegree) is $n-1$

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Directed subgraphs

A **subgraph** of $G=G(V,E)$ is a subset W of the vertex set V together with all of the arcs that connect pairs of vertices in W

Eg if $W=\{1,2,3\}$, the subgraph of



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Directed subgraph counts: the dyad census

The directed graph G has $n(n-1)/2$ subgraphs of size 2

Each subgraph of size 2 comprises a pair of vertices, and there are either 0, 1 or 2 arcs linking them:

subgraph

count



$N = \text{number of null dyads}$



$A = \text{number of asymmetric arcs}$



$M = \text{number of mutual arcs}$

Dyad census = (M, A, N) *count of the no. of each type of dyadic subgraph*

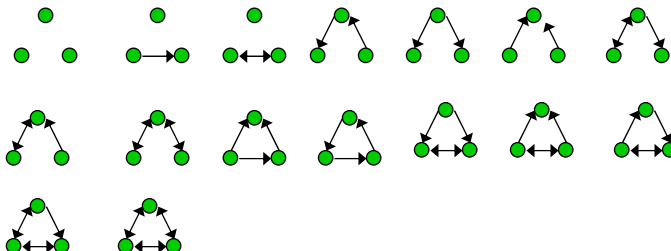
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Directed subgraph counts: the triad census

The graph G has $n(n-1)(n-2)/6$ subgraphs of size 3

Each subgraph of size 3 comprises a triple of vertices, and there are 16 possible forms:

subgraphs



Triad census: counts of each of the 16 forms across all subgraphs of G

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Paths and semipaths

A **path** from vertex a to vertex b is an ordered sequence

$$a=v_0, v_1, \dots, v_m=b$$

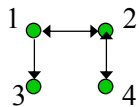
of distinct vertices in which each adjacent pair (v_{j-1}, v_j) is linked by an **arc**. The **length** of the path is m

A **semipath** from vertex a to vertex b is an ordered sequence

$$a=v_0, v_1, \dots, v_m=b$$

of distinct vertices in which either (v_{j-1}, v_j) and/or (v_j, v_{j-1}) is linked by an **arc**. The **length** of the semipath is m

e.g.



there is a path from 2 to 3 of length 2

there is no path from 3 to 4

there s a semipath from 3 to 4 of length 3

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Strong and weak connectedness; strong and weak components

If there is a path from vertex a to vertex b , b is **reachable** from a

If each vertex in G is reachable from each other vertex, then G is **strongly connected**

If there is a semipath from each vertex in G to each other vertex, then G is **weakly connected**

A **strong (weak) component** of G is a maximal strongly (weakly) connected subgraph (ie a strongly (weakly) connected subgraph with vertex set W for which no larger set Z containing W is strongly (weakly) connected)

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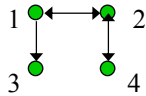
Geodesics

A *geodesic* from a to b is a path of minimum length

The *geodesic distance* d_{ab} between a and b is the length of the geodesic

If there is no path from a to b , the geodesic distance is *infinite*

For the directed graph



The geodesic distances are:

$$d_{12} = 1, d_{13} = 1, d_{14} = 2,$$

$$d_{21} = 2, d_{23} = 2, d_{24} = 1,$$

$$d_{31} = \text{infinite}, d_{32} = \text{infinite}, d_{34} = \text{infinite}, d_{41} = 2, d_{42} = 1, d_{43} = 3$$

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Centrality in directed graphs

As for graphs, but note that:

Indegree and *outdegree* centrality replace degree centrality

Eigenvector centrality is only computed by UCINET for graphs

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Some special directed graphs

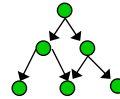
Empty and *complete* directed graphs



Cycle



Acyclic directed graph: a directed graph with no cycles



Transitive directed graph: every two path is accompanied by a direct path

