



Figure 1: Minimum spanning trees (question 1)

**Combinatorial Optimisation    HT 2010**  
**Problem set 1, MSc in Applied Statistics, for week 5**

- 1 Consider the graph shown in the figure, with nodes numbered 1 to 7 left to right along the top row, then along the next row and last the bottom node.
  - (a) Use Kruskal's algorithm to find a minimum spanning tree in the network illustrated in the figure; and then use Prim's algorithm to do the same, starting at node 1. In each case give the sequence of edges added to the tree.
  - (b) Now ignore the above minimum spanning trees. Use the green rule (once) and the red rule (twice) to show that there is a minimum spanning tree containing the edge 23 (cost 2) and not containing the edges 35 (cost 5) or 36 (cost 6).
  
- 2 The matrix  $A$  below gives the arc lengths in a network on the six nodes  $1, \dots, 6$  (where a dot indicates the absence of an arc); and similarly for the matrix  $B$ .
  - (a) For the matrix  $A$ , find the shortest path lengths from node 1 to each other node in the network, by Dijkstra's method.
  - (b) For the matrix  $B$ , find all shortest path lengths by using Floyd's method. [Why does the second row of the matrix never change? Why can there be no change in the iteration with  $k = 2$ ?]
  - (c) Now consider the matrix  $A$  again, and suppose that the length of arc 53 is changed from 2 to -6. Use Floyd's method to detect a negative cycle. [With  $k = 1$  there can be no changes, so (for hand calculation) we can jump to  $k = 2$ . You should detect a negative cycle with  $k = 4$ .]

$$A = \begin{pmatrix} \cdot & 3 & \cdot & 2 & \cdot & \cdot \\ \cdot & \cdot & 7 & 2 & \cdot & \cdot \\ \cdot & 1 & \cdot & 1 & \cdot & 3 \\ \cdot & 2 & \cdot & \cdot & 4 & \cdot \\ \cdot & \cdot & 2 & 3 & \cdot & 6 \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot \end{pmatrix} \quad B = \begin{pmatrix} 0 & 3 & -2 & 2 \\ \cdot & 0 & \cdot & \cdot \\ \cdot & 5 & 0 & 3 \\ 1 & 1 & \cdot & 0 \end{pmatrix}$$

- 3** If an organisation allocates  $i$  of its sales representatives to region  $j$  the profit from that region is given by  $a_{ij} \geq 0$  ( $1 \leq i \leq m$ ,  $1 \leq j \leq n$ ). Describe *very briefly* how dynamic programming gives an algorithm for finding the optimal allocation of sales representatives, which uses no more than  $f(m, n)$  integer comparisons and additions, where  $f$  is a polynomial which should be found.

Illustrate the algorithm by using it to find an optimum allocation when the  $a_{ij}$  are presented as the following array:

$i \setminus j$	1	2	3	4	5
0	0	0	0	0	0
1	21	39	50	37	42
2	39	60	110	53	63
3	73	79	139	72	89
4	90	105	180	120	130

- 4** Describe an algorithm for the problem  $P^\lambda$ , to maximise

$$\sum_{i=1}^n r_i(x_i, y_i) - \lambda \sum_{i=1}^n y_i$$

subject to  $\sum_{i=1}^n x_i = \alpha$ ,  $x_i \geq 0$ ,  $0 \leq y_i \leq \beta$ , and  $x_i$  and  $y_i$  integral for each  $i$ . How long will your method take? (Assume that you have a table for the values  $r_i(x, y)$ .)

- 5** Consider a flowshop with machines  $M_1, \dots, M_m$ .
- (i) Show that for a problem  $n/m/F/F_{max}$  there is an optimal schedule with the same job sequence on machines  $M_{m-1}$  and  $M_m$ .
- (ii) Show that for a problem  $n/3/F/F_{max}$  there is an optimal schedule with the same job sequence on all three machines.
- 6** Consider a problem  $n/1//\bar{T}$ . Given a set  $Q$  of jobs let  $H(Q)$  be the minimum value of the total tardiness for scheduling the jobs in  $Q$ . Give a dynamic programming recurrence for  $H(Q)$ . Use this approach to solve the following  $4/1//\bar{T}$  problem:

$i$	1	2	3	4
$p_i$	9	12	7	4
$d_i$	15	19	23	21